ON LOCAL LINEAR CONVERGENCE OF ELEMENTARY ALGORITHMS WITH SPARSITY CONSTRAINTS

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Central Problem: ℓ_0 minimization

Let $\mathcal{A}: \mathbb{R}^{m \times n} \to \mathbb{R}^r$ (not necessarily linear). Solve

(\mathcal{P}_0)		rank(x)
	subject to	$\mathcal{A}(x) = b$

where rank(x) := $\|\sigma(x)\|_0$:= $\sum_j \text{sign}(\sigma(x)_j)$ with sign (0) := 0 for $\sigma(x)$ the vector of singular values of x.

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Combinatorial optimization problem \implies NP-complete (Natarajan, 1995).

Central Problem: ℓ_0 minimization

Relaxations For 0 solve instead

where
$$\|x\|_{p} := \left(\sum_{j} \sigma(x)_{j}^{p}\right)^{1/p}$$
.
(For $p = 1$ and $m > 1$, $\|x\|_{1}$ is the *nuclear norm*.)

- ► m = 1, p = 1, A linear and satisfies the Restricted Isometry Property (RIP): optimal solutions of P₀ and P₁ coincide (Candés & Tao, 2006-2007).
- ► m = 1, 0 Isometry Property (RIP): optimal solutions of P₀ and P_p coincide (LaiWaing, 2010).

Central Problem: ℓ_0 minimization

Rewieghted/greedy algorithms

For m = 1, A linear, p = 1 and $\alpha > 0$ fixed, solve

$$(\mathcal{P}_{1,L})$$
 minimize $f(x,L)$
 $(x,L) \in \mathbb{R}^n \times \mathbb{R}^n$ subject to $Ax = b$

where

$$f(\boldsymbol{x}, \boldsymbol{L}) := \sum_{j=1}^{n} L_{j} |\boldsymbol{x}_{j}| + \iota_{\mathbb{R}_{+}}(\boldsymbol{L}) + \frac{1}{2} \|\boldsymbol{L} - \boldsymbol{\alpha}\|^{2}$$

- Orthogonal Matching Pursuit (OMP) (Chen, Donoho & Saunders, 1998-99)
- steepest subgradient descent algorithm with exact linesearch and dynamic weights applied to the *dual* (Borwein & L., 2011) (recovers OMP as a limiting case).
- see also Fornasier, Rauhut & Ward (2011), Candés, Wakin, & Boyd (2008), ...

Nonconvex Feasibility

find $x \in M_{\epsilon} \cap C_s$

where

$$M_{\epsilon} = \left\{ x \in \mathbb{R}^{m imes n} \mid d_{\mathbb{R}^{r}}(\mathcal{A}(x), b) \leq \epsilon
ight\}$$

for $d_{\mathbb{R}^r}$: $\mathbb{R}^r \times \mathbb{R}^r \to \mathbb{R}_+$ a distance function, and

$$\mathcal{C}_{\boldsymbol{s}} := \left\{ \boldsymbol{x} \in \mathbb{R}^{m \times n} \mid \operatorname{rank}(\boldsymbol{x}) \leq \boldsymbol{s} \right\}.$$

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- Combettes&Trussell (1990)
- Beck&Teboulle (2011).

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Central Algorithm: Method of Alternating Projections (MAP)

Consider more generally the problem

find $x \in A \cap B$

for two nonempty closed subsets A, B of a Hilbert space \mathcal{H} .

For a closed nonempty subset $C \subset \mathcal{H}$ define the *projection* of *x* onto *C* by

$$P_C: x \mapsto \operatorname{argmin}_{y \in C} \{ \|x - y\| \}$$

Method of Alternating Projections

given $x^0 \in \mathcal{H}$ generate the sequence $\{x^k\}_{k=0}^{\infty}$ in \mathcal{H} by

$$b^k \in P_B a^k$$
, $a^{2k} \in P_A b^{k-1}$.

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Method of Alternating Projections (MAP)

- ► If *A* and *B* are subspaces the MAP sequence converges strongly to $P_{A \cap B} x^0$ (von Neumann, 1933).
- If A and B are closed and convex and if one of the sets is compact the MAP sequence converges to a fixed point (Cheney-Goldstein 1959).
- ▶ If *A* and *B* are closed and convex with nonempty intersection the MAP sequence converges weakly to a point $\overline{x} \in A \cap B$ (Bregman, 1965).

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Method of Alternating Projections (MAP)

- ► If the angle between the subspaces *A* and *B* is positive then convergence is linear (Aronszajn, 1950).
- If A and B are closed and convex with int A ∩ B ≠ Ø the MAP sequence converges linearly to a point x̄ ∈ A ∩ B (Gubin-Polyak-Raik, 1967).
- If A and B are manifolds in E with T_A(x̄) + T_B(x̄) = E then for initial points x⁰ close to x̄ the MAP sequence converges linearly (Lewis & Mallick, 2008).
- If A ⊂ E is closed, B ⊂ E is superregular and
 N_A(x) ∩ −N_B(x) = {0} then for initial points x⁰ close to x the MAP sequence converges linearly (Lewis, L. & Mallick, 2009).

Method of Alternating Projections (MAP)

Local convergence of MAP requires two things:

- ► regularity of the intersection $A \cap B \implies$ constraint qualifications
- ▶ regularity of the sets *A* and *B* near the intersection.

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Both of these notions require the normal cone.

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Restricted normal cones (Bauschke, L., Phan & Wang, 2012)

Let *A* and *B* be nonempty subsets of a Euclidean space \mathbb{E} , and let *a* and *u* be in \mathbb{E} . If $a \in A$, we define the restricted normal cones of *A* at *a* as follows:

1. The B-restricted proximal normal cone of A at a is

$$\widehat{N}_{A}^{B}(a) := \operatorname{cone}\left(\left(B \cap P_{A}^{-1}a\right) - a\right) = \operatorname{cone}\left(\left(B - a\right) \cap \left(P_{A}^{-1}a - a\right)\right)$$

 The *B*-restricted normal cone N^B_A(a) is implicitly defined by u ∈ N^B_A(a) if and only if there exist sequences (a_n)_{n∈N} in A and (u_n)_{n∈N} in N^B_A(a_n) such that a_n → a and u_n → u.
 If a ∉ A, then all normal cones are defined to be empty.



Joint constraint qualification number (CQ-number)

Let A, \widetilde{A} , B, \widetilde{B} , be nonempty subsets of \mathbb{E} , let $c \in \mathbb{E}$, and let $\delta \in \mathbb{R}_{++}$. The *CQ-number* at c associated with $(A, \widetilde{A}, B, \widetilde{B})$ and δ is

$$egin{aligned} & heta_\delta := heta_\deltaig(m{A}, \widetilde{m{A}}, m{B}, \widetilde{m{B}}ig) \ &:= \sup\left\{ig\langle u, v
angle \; \left| egin{aligned} & u \in \widehat{N}_A^{\widetilde{B}}(m{a}), v \in -\widehat{N}_B^{\widetilde{A}}(m{b}), \|u\| \leq 1, \|v\| \leq 1, \ \|m{a} - m{c}\| \leq \delta, \|m{b} - m{c}\| \leq \delta. \end{aligned}
ight\} \end{aligned}
ight.$$

For nontrivial collections¹ $\mathcal{A} := (A_i)_{i \in I}, \widetilde{\mathcal{A}} := (\widetilde{A}_i)_{i \in I},$ $\mathcal{B} := (B_j)_{j \in J}, \widetilde{\mathcal{B}} := (\widetilde{B}_j)_{j \in J}$ of nonempty subsets of \mathbb{E} , the *joint-CQ-number* at $c \in \mathbb{E}$ associated with $(\mathcal{A}, \widetilde{\mathcal{A}}, \mathcal{B}, \widetilde{\mathcal{B}})$ and $\delta > 0$ is

$$\theta_{\delta} = \theta_{\delta} \big(\mathcal{A}, \widetilde{\mathcal{A}}, \mathcal{B}, \widetilde{\mathcal{B}} \big) := \sup_{(i,j) \in I \times J} \theta_{\delta} \big(\mathcal{A}_i, \widetilde{\mathcal{A}}_i, \mathcal{B}_j, \widetilde{\mathcal{B}}_j \big).$$

¹The collection $(A_i)_{i \in I}$ is said to be *nontrivial* if $I \neq \emptyset$.

The *Friedrichs angle* between subspaces *A* and *B* is the number in $[0, \frac{\pi}{2}]$ whose cosine is given by

$$c(A,B) :=$$

 $\sup \left\{ |\langle a,b \rangle | \ \left| \ a \in A \cap (A \cap B)^{\perp}, b \in B \cap (A \cap B)^{\perp}, \|a\|, \|b\| \le 1 \right\}
ight\}$

CQ-number of two (affine) subspaces and Friedrichs angle (Bauscke, L., Phan & Wang 2012)

Let *A* and *B* be affine subspaces of \mathbb{E} , and let $\delta > 0$. Then

$$\theta_{\delta}(A, A, B, B) = \theta_{\delta}(A, \mathbb{E}, B, B) = \theta_{\delta}(A, A, B, \mathbb{E}) = c(A, b) < 1.$$
(2)
Moreover, if A and B are affine subspaces of \mathbb{E} with $\overline{x} \in A \cap B$,
and $\delta > 0$, then (2) holds at \overline{x} .

Angle of intersection



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special case: two distinct lines through the origin

Suppose that w_a and w_b are two vectors in \mathbb{E} such that $||w_a|| = ||w_b|| = 1$. Let $A := \mathbb{R}w_a$, $B := \mathbb{R}w_b$, and $\delta > 0$. Assume that $A \cap B = \{0\}$. Then the CQ-number at 0 is

$$\theta_{\delta}(\boldsymbol{A},\boldsymbol{A},\boldsymbol{B},\boldsymbol{B}) = \theta_{\delta}(\boldsymbol{A},\mathbb{E},\boldsymbol{B},\boldsymbol{B}) = \theta_{\delta}(\boldsymbol{A},\boldsymbol{A},\boldsymbol{B},\mathbb{E}) = |\langle \boldsymbol{w}_{\boldsymbol{a}},\boldsymbol{w}_{\boldsymbol{b}}\rangle| < 1.$$

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Superregularity (Lewis, L. & Mallick, 2009) We say that *C* is (A, ϵ, δ) -regular at $c \in \mathbb{E}$ if $\epsilon \ge 0, \delta > 0$, and

$$\begin{array}{c} (y,b) \in B \times B, \\ \|y-c\| \leq \delta, \|b-c\| \leq \delta, \\ u \in \widehat{N}_{B}^{\mathcal{A}}(b) \end{array} \right\} \quad \Rightarrow \quad \langle u, y-b \rangle \leq \epsilon \|u\| \cdot \|y-b\|.$$

$$(3)$$

If *B* is $(\mathbb{E}, \epsilon, \delta)$ -regular at *c*, then we also simply speak of (ϵ, δ) -regularity.

Example: sets for which the projector is single-valued on neighborhoods of the set (prox-regular sets) are (ϵ, δ) -regular. Convex sets are $(0, \infty)$ -regular.

Harvest Time

Linear convergence of MAP

Let $\mathcal{A} := (A_i)_{i \in I}$ and $\mathcal{B} := (B_j)_{j \in J}$ be nontrivial collections of nonempty closed subsets of \mathbb{E} , and let $\widetilde{\mathcal{A}} := (\widetilde{A}_i)_{i \in I}$ and $\widetilde{\mathcal{B}} := (\widetilde{B}_j)_{j \in J}$ be reasonable collections of subsets of \mathbb{E} , $\widetilde{\mathcal{A}} := \bigcup_{i \in I} \widetilde{\mathcal{A}}_i$ and $\widetilde{\mathcal{B}} := \bigcup_{j \in J} \widetilde{\mathcal{B}}_j$. Let

 $c \in (\cup_{i \in I}(A_i)) \cap (\cup_{j \in J}(B_j))$

and assume that there exists $\delta > 0$ such that

- **1.** \mathcal{A} is $(\tilde{B}, 0, 3\delta)$ -joint-regular at *c*;
- **2.** \mathcal{B} is $(\widetilde{A}, 0, 3\delta)$ -joint-regular at c; and
- **3.** $\theta < 1$, where $\theta := \theta_{3\delta}$ is the joint-CQ-number at *c* associated with $(\mathcal{A}, \widetilde{\mathcal{A}}, \mathcal{B}, \widetilde{\mathcal{B}})$

Harvest Time

Linear convergence of MAP

Suppose also that the starting point of the MAP b_{-1} satisfies $||b_{-1} - c|| \le \frac{(1-\theta)\delta}{6(2-\theta)}$. Then $(a_n)_{n\in\mathbb{N}}$ and $(b_n)_{n\in\mathbb{N}}$ converge linearly to some point in $\bar{c} \in A \cap B$ with $||\bar{c} - c|| \le \delta$ and rate θ^2 ; in fact,

$$(\forall n \geq 1) \quad \max \left\{ \|a_n - \bar{c}\|, \|b_n - \bar{c}\| \right\} \leq \frac{\delta}{2 - \theta} (\theta^2)^{n-1}$$

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Normal cone to C_s

Recall that $C_s := \{x \in \mathbb{R}^{m \times n} | \operatorname{rank}(x) \le s\}$ and define $\operatorname{Supp}(x) := \operatorname{range}(x^T)$.

Normal cone to C_s (L. 2012)

Let
$$r := \min\{m, n\}$$
, fix $s \in \{0, 1, ..., r\}$ and define the set
 $C_s := \{x \in \mathbb{R}^{m \times n} \mid \operatorname{rank}(x) \le s\}$. At a point $\overline{x} \in C_s$
 $N_{C_s}(\overline{x}) = \{v \in \mathbb{R}^{m \times n} \mid v \in \operatorname{Supp}(\overline{x})^{\perp} \text{ and } \operatorname{rank}(v) \le r - s\}$

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Moreover, $N_{\mathcal{C}_s}(\overline{x}) = \widehat{N}_{\mathcal{C}_s}(\overline{x})$ at every \overline{x} with $\operatorname{rank}(\overline{x}) = s$.

Normal cone to C_s : vector case, $x \in \mathbb{R}^n$

Define

$$\mathcal{J}:=2^{\{1,2,\dots,n\}}$$
 and $\mathcal{J}_{s}:=\mathcal{J}(s):=\{J\in\mathcal{J}\mid |J|=s\}$

and set

$$(\forall J \in \mathcal{J}) \quad \mathcal{C}_J := \operatorname{span}\left\{ e_j \mid j \in J \right\}.$$

Define the collections

$$\mathcal{C} := \widetilde{\mathcal{C}} := (\mathcal{C}_J)_{J \in \mathcal{J}_s}$$
 and $\mathcal{M} := \widetilde{\mathcal{M}} := M := \{x \in \mathbb{R}^n \mid Ax = b\}.$
Clearly,

$$\mathcal{C}_{\boldsymbol{s}} := \left\{ \boldsymbol{x} \in \mathbb{R}^n \mid \|\boldsymbol{x}\|_0 \leq \boldsymbol{s} \right\} = \widetilde{\mathcal{C}}_{\boldsymbol{s}} := \bigcup_{J \in \mathcal{J}_{\boldsymbol{s}}} \boldsymbol{C}_J$$

and

$$M=\left\{x\in\mathbb{R}^n\mid Ax=b\right\}.$$

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Normal cone to C_s : vector case, $x \in \mathbb{R}^n$

Then

$$\begin{split} \mathcal{N}_{\mathcal{C}_s}(\overline{x}) &= \left\{ v \in \mathbb{R}^n \ \Big| \ \|v\|_0 \leq n-s \text{ and } v \in \operatorname{Supp}{(\overline{x})^{\perp}} \right\} \\ &= \bigcup_{I(\overline{x}) \subseteq J \in \mathcal{J}_s} C_J^{\perp}. \end{split}$$

Consequently, if $\|\overline{x}\|_0 = s$, then $N_{\mathcal{C}_s}(\overline{x}) = \operatorname{Supp}(\overline{x})^{\perp} = C_{l(\overline{x})}^{\perp}$.

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Application to sparse optimization

MAP for sparse-affine feasibility converges locally linearly

Let C_s , C, \widetilde{C} M, \mathcal{M} and $\widetilde{\mathcal{M}}$ be defined as above. Suppose that $s \leq n-1$, that $\overline{x} \in C_s \cap M$, and fix $\delta \in (0, \overline{\delta})$ for $\overline{\delta} := \frac{1}{3} \min \{ d_{C_J}(\overline{x}) \mid \overline{x} \notin C_J, J \in \mathcal{J}_s \}$. Then

 $\overline{\delta} = \frac{1}{3} \min \left\{ |\overline{x}_j| \mid j \in I(\overline{x}) \right\}$

and

$$\theta_{3\delta}(\mathcal{C},\widetilde{\mathcal{C}},M,M) = \max \{ c(C_J,M) \mid \overline{x} \in C_J, J \in \mathcal{J}_s \} < 1,$$

where $\theta_{3\delta}$ denotes the joint-CQ-number at \overline{x} associated with $(\mathcal{C}, \widetilde{\mathcal{C}}, M, M)$.

Application to sparse optimization

MAP for sparse-affine feasibility converges locally linearly

Generate the sequences $\{c_k\}_{k\in\mathbb{N}}$ and $\{m_k\}_{k\in\mathbb{N}}$ in \mathbb{R}^n by the MAP algorithm

$$c_k = P_{\mathcal{C}_s}(m_{k-1})$$
 and $m_k = P_M(c_k)$

where

$$M = \left\{ x \in \mathbb{R}^{m imes n} \mid Ax = b
ight\} ext{ and } \mathcal{C}_{s} := \left\{ x \in \mathbb{R}^{n} \mid \|x\|_{0} \leq s
ight\}.$$

Suppose the starting point of the MAP $m_{-1} \in M$ satisfies $||m_{-1} - \overline{x}|| \leq \frac{(1-\overline{\theta})\delta}{6(2-\overline{\theta})}$. Then $(c_k)_{k\in\mathbb{N}}$ and $(m_k)_{k\in\mathbb{N}}$ converge linearly to some point in $\overline{c} \in C_s \cap M \cap \mathbb{B}(\overline{x}, \delta)$ with rate $\overline{\theta}^2$.

MAP for sparse-affine feasibility converges locally linearly

Remarks

- Regularity of the intersection is not an assumption of the theorem; it is *automatically* satisfied. This is in contrast to the results of Lewis&Mallick (2008) and Lewis,L.&Mallick (2009) where the required regularity is assumed to hold. Simple examples illustrate that the notions of regularity developed in those works are not satisfied
- 2. Our analysis is the first (to our knowledge) to yield a quantification of the neighborhood on which local linear convergence is guaranteed.
- Finding the local neighborhood on which linear convergence is guaranteed may well be tantamount of finding the sparsest solution; however, it does open the door to justify combining the MAP with more aggressive algorithms such as Douglas-Rachford

Thanks for your attention.





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